

# Dual-Zone Hard-Core Model for RTS/CTS Handshake Analysis in WLANs

Yi Zhong, *Senior Member, IEEE*, Zhuoling Chen, Wenyi Zhang, *Senior Member, IEEE*, Martin Haenggi, *Fellow, IEEE*

**Abstract**—This paper introduces a new stochastic geometry-based model to analyze the Request-to-Send/Clear-to-Send (RTS/CTS) handshake mechanism in wireless local area networks (WLANs). We develop an advanced hard-core point process model, termed the dual-zone hard-core process (DZHCP), which extends traditional hard-core models to capture the spatial interactions and exclusion effects introduced by the RTS/CTS mechanism. This model integrates key parameters accounting for the thinning effects imposed by RTS/CTS, enabling a refined characterization of active transmitters in the network. Analytical expressions are derived for the intensity of the DZHCP, the mean interference, and an approximation of the success probability, providing insight into how network performance depends on critical design parameters. Our results demonstrate that the Type II RTS/CTS mechanism significantly reduces mean interference by introducing additional protection regions, effectively mitigating the impact of nearby interferers. Compared to CSMA-based schemes, it achieves a lower mean interference level while maintaining a comparable or higher active node density.

## I. INTRODUCTION

### A. Motivations

As a robust complement to the 5G and future 6G mobile communication networks [2], wireless local area networks (WLANs) will continue to play a pivotal role in data transmission processes. The widespread adoption of WLANs in recent years has significantly improved personal access to wireless networks. The IEEE 802.11 media access control (MAC) standard [3] introduces both physical and virtual carrier sensing techniques to prevent data collisions, a critical issue in densely populated network environments. Physical carrier sensing allows a transmitter to detect active transmissions within its range and to defer its activity if necessary [4]. In contrast, virtual carrier sensing, implemented through the Request-to-Send/Clear-to-Send (RTS/CTS) handshake mechanism, addresses the hidden terminal problem, where a node visible to the receiver is obscured from the transmitter [5]. This mechanism establishes a protective region around both the transmitter and receiver, effectively minimizing collision risks in overlapping transmission regions (Figure 1). RTS/CTS effectively prevents packet collisions by reserving the channel

Yi Zhong and Zhuoling Chen are with the School of Electron. Inf. & Commun., Huazhong University of Science and Technology, Wuhan, China. Wenyi Zhang is with the University of Science and Technology of China. Martin Haenggi is with the University of Notre Dame, US. The research has been supported by the National Natural Science Foundation of China (NSFC) grant No. 62471193 and No. 62231022, and the Wuhan Science and Technology Program (No. 2024030803010178). The computation is completed in the HPC Platform of Huazhong University of Science and Technology. Preliminary findings of this work were presented in the IEEE International Conference on Communications (ICC) [1].

The corresponding author is Yi Zhong (yzhong@hust.edu.cn).

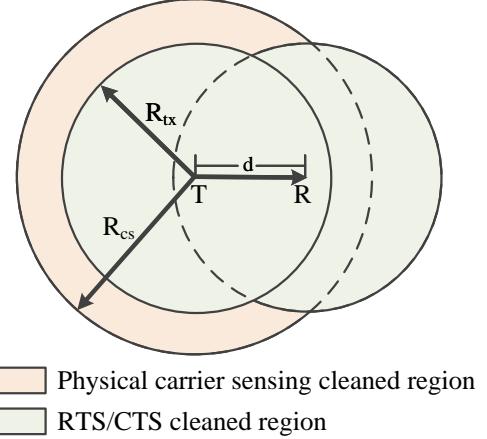


Fig. 1. Illustration of a transceiver pair employing the RTS/CTS mechanism, showing the dual-zone exclusion regions for interference management. The transmitter (T) and receiver (R) are separated by a distance  $d$ . The orange region, with radius  $R_{cs}$ , corresponds to the physical carrier sensing range, which prevents nearby transmitters from causing interference at the transmitter (T). The green region, with radius  $R_{tx}$  and centered at the receiver (R), denotes the RTS/CTS virtual carrier sensing region, which protects the receiver from potential interference by blocking transmissions within this zone.

for data transmission, particularly in dense networks. While this mechanism adds control overhead, its benefits in reducing retransmissions and improving overall network efficiency often outweigh the costs, especially for larger data packets.

Despite considerable research on the performance of 802.11 WLANs [6]–[10], a precise analysis of the RTS/CTS mechanism’s network impact remains challenging. The complexity arises from the randomness in the spatial distribution of nodes and the significant correlations between active transmitters. Understanding the impact of RTS/CTS design parameters, such as carrier sensing ranges on network interference, is crucial for optimizing network performance.

This paper introduces a novel spatial distribution model specifically designed for analyzing the RTS/CTS mechanism. Traditional models, such as the homogeneous Poisson point process (PPP) and the Matérn hard-core process (MHCP) [11]–[13], fall short due to the unique spatial dependencies introduced by RTS/CTS. Under RTS/CTS mechanism, communication behaviors between devices introduce spatial dependencies, causing node distributions to deviate from the assumption of complete randomness. However, PPP models are only suitable for simulating network scenarios with independently distributed nodes and are inadequate to accurately capture the spatial correlations introduced by RTS/CTS mechanism. Moreover, MHCP models, while accounting for

spatial exclusion effects between nodes, are primarily designed to establish protective zones around transmitters to reduce potential interference. In contrast, the RTS/CTS mechanism sets up protective zones for both the transmitter and the receiver. Our approach utilizes a marked point process that more accurately characterizes the spatial relationships between transmitters and receivers within this framework. Each node in this process not only indicates the position of a transmitter but also includes a mark representing the receiver's relative location. The influence of RTS/CTS on this model is analogous to the effect observed in Matérn hard-core models on PPPs [14]. Using the proposed model, we derive formulas for calculating the mean interference, accounting for all potential network configurations. Moreover, we employ approximations to derive the success probability of transmissions. Although the asymptotic gain approximation method has been proved to be effective in cellular networks, prior research has primarily focused on nearest access scenarios [15]–[18]. To the best of our knowledge, this study is the first to extend this method to the bipolar dual hard-core process. This extension provides a novel approach to efficiently approximating the success probability in such settings, thereby offering new insights for the modeling and optimization of complex communication systems.

### B. Related Work

The effective management of network interference and optimization of access resources in wireless networks has been a focal area of research [19]–[24]. In [19], the authors investigated the influence of the RTS threshold on the network sum rate, under the condition that the network always operates in the access mode. The authors of [20] propose a new chain RTS/CTS scheme for multihop communication in aerial networks, which achieves 37% throughput improvement in multihop links with 68% increased connection establishment time compared to existing IEEE 802.11 Independent Basic Service Set (IBSS). In [22], the authors proposed an optimal power control strategy designed to maximize the access probability in carrier sense multiple access (CSMA) protocols by utilizing a max-interference model. While this approach marks a significant advancement in power control tactics, it notably omits a comprehensive statistical analysis of interference, which could otherwise provide deeper insights into the dynamic interplay between power levels and interference patterns across the network. The authors of [23] derive a closed-form interference avoidance criterion based on a multi-frequency RTS/CTS scheme, overlooking the spatial correlation among transmitters [24].

Moreover, traditional interference models for CSMA protocols, such as those based on Matérn's hard-core process discussed in [13], [25], [26], have been extensively used to analyze network behaviors. These models are appreciated for their theoretical simplicity and tractability. Particularly in [25], the outage probability was derived using an approximation to the PPP, offering a manageable yet somewhat imprecise methodology for understanding network resilience. Conversely, [13] advanced this by deriving more accurate

expressions for the mean interference within the Matérn hard-core framework, enhancing the model's utility in practical applications. The process has been utilized as a fundamental model to emulate a variety of spatial configurations and scenarios, particularly where entities exhibit a minimum exclusion distance from one another [27]–[30]. But the intensity underestimation flaw of Matérn's hard-core process Type II has not been addressed in the literature, mentioned in [31] and [32]. In [33], the authors introduced a modified hard-core point process inspired by the definition of Matérn's hard-core process Type III to mitigate the flaw. Building upon the novel framework presented [33], expressions for outage probability and throughput were derived at the user pair served using NOMA [34]. The modified Matérn hard-core point process Type III model, which is an improved simulation method based on Matérn hard-core point process Type III, was discussed in order to resolve the inaccurately reflect of modified hard-core point process in [35].

However, both approaches encounter significant limitations when applied to analyzing networks utilizing the RTS/CTS mechanism. This mechanism introduces additional complexity due to the spatial dependencies among transmitters and their corresponding receivers. In such configurations, the receivers' locations are marked at transmitters and become integral in suppressing potential interference from neighboring transmissions. This intricacy renders traditional tools like the reduced second moment measure and the Ripley K-function, which [13] employed, less effective. These methods, while robust under simpler point process assumptions, struggle to capture the impact of RTS/CTS protocols on network performance due to their inability to fully integrate the spatial correlation introduced by RTS/CTS interactions. Recognizing these gaps, our work aims to extend the existing methodologies by incorporating a marked point process approach that explicitly considers the spatial configuration of both transmitters and receivers influenced by RTS/CTS mechanisms.

### C. Contributions

This study proposes an analytical framework for wireless networks based on the RTS/CTS handshake mechanism. Addressing identified gaps in existing research, we introduce new tools and methodologies that enhance the analysis and optimization of WLAN performance. The key findings and contributions of our study are summarized as follows:

- We introduce the dual-zone hard-core process (DZHCP) as a new stochastic model developed for networks utilizing RTS/CTS protocols. This model distinctively captures dual protection regions around transmitters and receivers, enabling a fine-grained analysis of spatial interactions and interference dynamics.
- We have enhanced classical statistical measures, such as the reduced second moment measure and Ripley's K-function, to better capture the characteristics of marked point process models in our analysis. These improvements allow for a more accurate representation of spatial interactions governed by the RTS/CTS mechanism. Building on these refined models and measures, we derive exact expressions for the mean interference.

- We further contribute by deriving an approximation for the success probability of wireless links under the RTS/CTS mechanism. This is the first to generalize the asymptotic gain approximation method to ad hoc networks, providing an efficient approximation of success probabilities in bipolar models.

The remainder of this paper is organized as follows: Section II outlines our marked point process models for nodes under the RTS/CTS protocols, establishing the theoretical framework. Section III presents our main analytical results, including the intensity of the marked point process and the mean interference. Section IV presents our approximation for the success probability for the proposed DZHCP. Section V provides numerical illustrations and demonstrates their practical applications in WLAN environments.

## II. STATISTICAL MODEL FOR RTS/CTS NETWORKS

### A. Model Assumptions

We consider a densely populated geographic region with potential transmitter-receiver pairs operating within a shared frequency band, specifically focusing on a single subchannel of an 802.11 WLAN. Subchannels are assumed to be randomly and independently allocated, making this model applicable to multi-channel systems.

Each transmitter broadcasts at a common power level,  $P_t$ . Channel fading follows a Rayleigh model, with fading coefficients that remain constant within each time slot and are independent across time and space, modeled as exponentially distributed with unit mean. The fading between a transmitter at point  $x$  and a receiver at point  $z$  is represented by  $h_{xz}$ . For path loss, we use a deterministic function  $l(\cdot)$ , commonly  $l(r) = Ar^{-\alpha}$  where  $\alpha > 2$  is the path-loss exponent. Under this model, the received power at the receiver at  $z$  from a transmitter at  $x$  is  $h_{xz}l(|x - z|)$ .

### B. Dual-Zone Hard-Core Process Model

We model the locations of transmitters and receivers using a Poisson bipolar model [36], where each point in a PPP represents a potential transmitter, paired with a receiver located at a fixed distance  $d$  and a random angle  $\theta$  (see Figure 2). Although  $d$  is fixed in this basic model, it can be extended by allowing  $d$  to vary as a random variable. This assumption simplifies the analysis by ensuring uniform protection regions for all transceiver pairs, making it easier to derive key network performance metrics. In practical wireless networks, transmitter-receiver distances are often variable rather than fixed. Incorporating randomized link distances could provide a more realistic representation of network dynamics. Moreover, our model assumes omnidirectional RTS/CTS transmission, leading to circular protection regions. While this is a common simplification, many modern networks leverage beamforming, which creates directional transmission patterns and affects spatial interference characteristics. A promising direction for future research is to extend the model to account for beamforming and investigate its impact on interference management and network performance.

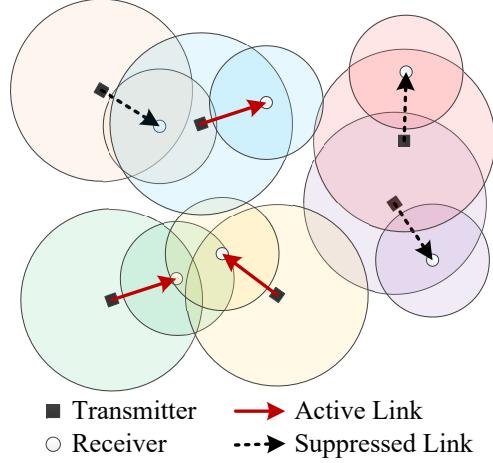


Fig. 2. Illustration of a bipolar model with RTS/CTS signaling and Type I thinning. In this model, each transmitter (represented as a black square) is paired with a corresponding receiver (depicted as a white circle). Active links are shown with solid red arrows, representing the connections that remain functional after the RTS/CTS thinning process. Suppressed links, depicted by dashed black arrows, represent connections that are blocked or “thinned” to prevent interference with neighboring active links, in accordance with the RTS/CTS mechanism.

In the RTS/CTS mechanism, when a transmitter is ready to send data, it first transmits an RTS frame to the intended receiver. If the receiver is available to accept data, it responds with a CTS frame as an acknowledgment. This process establishes a temporary protected region around both the transmitter and the receiver, preventing nearby nodes from initiating transmissions that could cause interference. Unlike traditional CSMA, where the protected region primarily exists around the transmitter, the RTS/CTS protocol extends this protection to both communication parties.

Another critical design parameter of the RTS/CTS mechanism is the RTS threshold, which determines whether the RTS/CTS handshake is used based on packet size. If the packet length exceeds a predefined threshold, the RTS/CTS mechanism is triggered; otherwise, the transmission proceeds without it. Setting the threshold too low results in excessive RTS/CTS overhead, reducing network efficiency due to additional control message exchanges. Conversely, a threshold that is too high may prevent RTS/CTS from being applied to large packets, increasing the risk of collisions as other devices may fail to sense ongoing transmissions.

Incorporating the RTS/CTS mechanism, each transceiver pair establishes an exclusion region consisting of two regions: a physical carrier sensing region and a virtual carrier sensing region. The physical region is a disk centered on the transmitter with radius  $R_{cs}$ , while the virtual region includes overlapping disks centered on both the transmitter and receiver, each with radius  $R_{tx}$ , where  $R_{tx} < R_{cs}$ . Virtual carrier sensing relies on the RTS/CTS handshake mechanism, where a node must successfully decode the RTS/CTS frames to interpret the transmission duration. In contrast, physical carrier sensing is based on detecting the received signal power, requiring only that the signal strength exceeds a threshold. Since decoding RTS/CTS frames demands a higher SNR than mere power

detection, the effective range of virtual carrier sensing is generally smaller than that of physical carrier sensing.

This dual-zone configuration simulates the protective effect of the RTS/CTS handshake mechanism, using the average received power (ignoring fading) to simplify exclusion region calculations. This approach parallels assumptions used in modeling the CSMA protocol with Matérn hard-core processes. Figure 1 illustrates these overlapping regions, highlighting the spatial dynamics governed by the RTS/CTS protocol within our model.

In our network model, we capture the effect of the RTS/CTS handshake mechanism through a process we term *RTS/CTS thinning*. This process selectively filters points from the Poisson bipolar process of potential transceivers, following specific RTS/CTS rules to adapt the network layout and reduce interference. We introduce two types of RTS/CTS thinning, inspired by the Matérn hard-core model but customized to the specifics of wireless protocols using RTS/CTS.

In *Type I thinning*, a transceiver pair is retained only if no other potential transmitter exists within its exclusion region. This region is defined by the combined physical and virtual carrier sensing regions. Type I thinning models a scenario similar to slotted-time protocols, where each time slot is isolated, and transceivers compete for access within each slot. If a potential collision is detected (i.e., another transmitter is found within the exclusion region), the transceiver pair's transmission is deferred to avoid interference in the current slot.

*Type II thinning* introduces a more dynamic approach by assigning each transceiver pair a random mark, typically viewed as a timestamp. A transceiver pair is retained for transmission only if no other potential transmitter with an earlier timestamp (i.e., a lower mark) is within its exclusion region. This approach more closely represents real-world transmission patterns, where transmissions depend on the relative timing of signals rather than strict time slots. Type II thinning thus manages interference by prioritizing established transmissions over newer attempts, creating a transmission landscape that dynamically adjusts to timing variations.

The distinction between these thinning types reflects different strategies for managing network traffic and interference. *Type I thinning* provides a straightforward, conservative approach to minimizing signal interference, which can come at the cost of reduced network throughput. By contrast, *Type II thinning* introduces greater flexibility by prioritizing transceivers based on their initiation times, potentially enhancing network efficiency. In practical applications, Type II thinning aligns more closely with the behavior of modern RTS/CTS protocols, balancing robust interference management with more efficient network utilization.

The proposed dual-zone hard-core processes share some similarities with Matérn hard-core processes—particularly in their generation and thinning methods—but also introduce key differences in the configuration of exclusion regions around each transceiver pair. In a Type I Matérn hard-core process, exclusion regions are regular circles with a fixed radius for all transmitters. This symmetry means that if one transmitter falls within the exclusion region of another, the reverse is also

true, causing both transmitters to either be retained or removed together.

In contrast, the exclusion regions in a Type I dual-zone hard-core process are irregular, resulting in asymmetrical interactions. If one transmitter lies within the exclusion region of another, it does not necessarily mean that the second transmitter will be within the exclusion region of the first. This irregularity allows for scenarios where, under Type I thinning in the dual-zone hard-core process, one transceiver pair may be suppressed while the other remains active.

The dual-zone hard-core model is well-suited for full-duplex communication systems, where simultaneous transmission and reception on the same frequency channel require effective interference management [37]. Unlike traditional hard-core models that enforce protection only around transmitters, the dual-zone model introduces exclusion regions for both transmitters and receivers, better capturing full-duplex interference dynamics.

### C. Mathematical Description

1) *Type I Dual-Zone Hard-Core Process*: The Type I dual-zone hard-core process leverages a dependently marked PPP, denoted as  $\tilde{\Phi}_{\text{PPP}}$ , to represent the spatial distribution and interaction rules for transceivers operating under the RTS/CTS mechanism. This process, defined over the Euclidean plane  $\mathbb{R}^2$ , uses an intensity  $\lambda_p$  to capture the density of potential transmitters and is structured as follows:

- Transmitter locations  $\Phi_{\text{PPP}}$ : Represents the set of potential transmitters in the network modeled by a PPP, denoted by  $\Phi_{\text{PPP}} = \{x_i\}$ .
- Receiver orientation  $\theta_i$ : Each transmitter located at  $x_i$  has an associated receiver placed at a fixed distance  $d$  in a randomly chosen direction. The orientation of each receiver relative to its transmitter is denoted by  $\theta_i$ , which is uniformly distributed in  $[0, 2\pi]$ . Together, the pair  $(x_i, \theta_i)$  uniquely determines each receiver's location.
- Medium access indicator  $e_i$ : Each transmitter-receiver pair  $(x_i, \theta_i)$  is associated with an access indicator  $e_i$ , which determines whether the pair is allowed to transmit, based on local network conditions.

For a given transmitter-receiver pair  $(x, \theta, e)$  in  $\tilde{\Phi}_{\text{PPP}}$ , we define its neighborhood  $\mathcal{N}(x, \theta, e)$  as

$$\begin{aligned} \mathcal{N}(x, \theta, e) &\triangleq \{(x', \theta', e') \in \tilde{\Phi}_{\text{PPP}} \setminus \{(x, \theta, e)\} : \\ &x' \in B_x(R_{\text{cs}}) \cup B_{x+(d \cos \theta, d \sin \theta)}(R_{\text{tx}})\}, \end{aligned} \quad (1)$$

where  $B_x(r)$  denotes a disk of radius  $r$  centered at  $x$ . This neighborhood function  $\mathcal{N}$  includes all potential transmitters that fall within the exclusion regions of the transmitter-receiver pair  $(x, \theta, e)$ , defined by the physical carrier sensing radius  $R_{\text{cs}}$  and the virtual carrier sensing radius  $R_{\text{tx}}$ .

The medium access indicator  $e$  for a transmitter-receiver pair  $(x, \theta)$  is then determined as

$$e \triangleq \mathbb{1}(\#\mathcal{N}(x, \theta, e) = 0), \quad (2)$$

where  $\#$  denotes the count of elements in the neighborhood. This indicator ensures that a transmitter-receiver pair is allowed to transmit only if no other transmitters are present

within its combined exclusion region. In other words,  $e = 1$  if the exclusion region is empty, allowing transmission, and  $e = 0$  otherwise.

The collection of all transceiver pairs that successfully meet the conditions of the RTS/CTS thinning process is defined as

$$\tilde{\Phi} \triangleq \{(x, \theta) : (x, \theta, e) \in \tilde{\Phi}_{\text{PPP}}, e = 1\}, \quad (3)$$

where only pairs with  $e = 1$  are retained, indicating they have passed the medium access check. The final set of active transmitters, denoted by  $\Phi$ , is

$$\Phi \triangleq \{x : (x, \theta) \in \tilde{\Phi}\}. \quad (4)$$

This filtered set  $\Phi$  represents the transmitters that have successfully passed the RTS/CTS thinning protocol. By restricting the set to these approved transmitters, the model effectively reduces potential interference within the network, allowing for more efficient and reliable communications.

2) *Type II Dual-Zone Hard-Core Process*: In contrast to Type I, the Type II dual-zone hard-core process incorporates the asynchronous nature of transmission attempts, modeling a more realistic scenario where the RTS/CTS mechanism operates based on chronological transmission order. This process uses a dependently marked PPP, denoted as  $\tilde{\Phi}_{\text{PPP}}$ , over  $\mathbb{R}^2$  with an intensity  $\lambda_p$ . The components of this process are as follows:

- Transmitter locations  $\Phi_{\text{PPP}}$ : The set of potential transmitters within the network modeled by a PPP, represented by  $\Phi_{\text{PPP}} = \{x_i\}$ .
- Receiver orientation  $\theta_i$ : The orientation of each receiver relative to its transmitter at location  $x_i$ , independently and identically distributed (i.i.d.) over  $[0, 2\pi]$ . Combined with a fixed transmitter-receiver distance  $d$ , this orientation determines the exact location of each receiver.
- Transmission initiation time marks  $\{m_i\}$ : A set of i.i.d. time marks, uniformly distributed over  $[0, 1]$ , representing the relative initiation times of transmissions for each transmitter. These marks are crucial for resolving access contention among transceivers.
- Medium access indicator,  $e_i$ : Indicates whether a transceiver pair is eligible for transmission based on local network conditions and the timing indicated by  $m_i$ .

For each node  $(x, \theta, m, e)$  in  $\tilde{\Phi}_{\text{PPP}}$ , the neighboring set  $\mathcal{N}(x, \theta, m, e)$  is defined as

$$\begin{aligned} \mathcal{N}(x, \theta, m, e) \triangleq & \{(x', \theta', m', e') \in \tilde{\Phi}_{\text{PPP}} \setminus \{(x, \theta, m, e)\} : \\ & x' \in B_x(R_{\text{cs}}) \cup B_{x+(d \cos \theta, d \sin \theta)}(R_{\text{tx}})\}, \end{aligned} \quad (5)$$

where  $B_x(r)$  denotes a disk centered at  $x$  with radius  $r$ . This defines the exclusion region for each node, which is influenced by both the physical carrier sensing distance  $R_{\text{cs}}$  and the virtual carrier sensing distance  $R_{\text{tx}}$ .

The eligibility of a node  $(x, \theta, m, e)$  to transmit, indicated by  $e$ , is

$$e \triangleq \mathbb{1} (\forall (x', \theta', m', e') \in \mathcal{N}(x, \theta, m, e), m < m'). \quad (6)$$

This condition ensures that a node can transmit only if its time mark  $m$  is earlier than any other potential transmitters within its exclusion regions.

Following the application of the Type II thinning based on these criteria, the set of active transceiver pairs,  $\tilde{\Phi}$ , is

$$\tilde{\Phi} \triangleq \{(x, \theta) : (x, \theta, m, e) \in \tilde{\Phi}_{\text{PPP}}, e = 1\}. \quad (7)$$

This filtered set  $\tilde{\Phi}$  represents transceiver pairs that successfully passed the Type II thinning criteria.

Finally, the collection of active transmitters that survive the thinning process is

$$\Phi \triangleq \{x : (x, \theta) \in \tilde{\Phi}\}. \quad (8)$$

This process effectively captures the asynchronous and competitive dynamics of network access managed by the RTS/CTS mechanism, providing a more realistic, time-sensitive model of network behavior.

The proposed dual-zone hard-core process assumes circular protection regions around both the transmitter and receiver. This approximation is suitable for transmissions without beamforming but may not accurately capture the interference regions in beamforming-enabled networks. Additionally, the Type II model inherits the limitation of the Matérn hard-core process of Type II, leading to a potential underestimation of the density of active transmitters. Future work may focus on refining the model to solve these issues.

### III. INTENSITY AND MEAN INTERFERENCE

In this section, we establish the node intensity and the mean interference experienced by a receiver. We introduce several key quantities that are crucial for our calculations:

- Exclusion region area  $V_o$ : This is the area of the exclusion region around a transceiver pair, which is essential for calculating the density of active transmitters:

$$V_o = (\pi - \xi_1)R_{\text{cs}}^2 + (\pi - \xi_2)R_{\text{tx}}^2 + dR_{\text{cs}} \sin \xi_1, \quad (9)$$

where  $\xi_1 = \arccos\left(\frac{d^2 + R_{\text{cs}}^2 - R_{\text{tx}}^2}{2dR_{\text{cs}}}\right)$  and  $\xi_2 = \arccos\left(\frac{d^2 + R_{\text{tx}}^2 - R_{\text{cs}}^2}{2dR_{\text{tx}}}\right)$ , as illustrated in Figure 1.

- Spatial criteria for interference  $S_1$ ,  $S_2$ , and  $S_3$ : Define the spatial criteria for transmitters within the exclusion regions, helping to determine interference conditions:

$$\begin{aligned} - S_1 &= \{(r, \beta, \theta) : r \leq R_{\text{cs}}\}, \\ - S_2 &= \{(r, \beta, \theta) : r^2 - 2rd \cos \beta + d^2 \leq R_{\text{tx}}^2\}, \\ - S_3 &= \{(r, \beta, \theta) : r^2 + 2rd \cos(\beta - \theta) + d^2 \leq R_{\text{tx}}^2\}. \end{aligned}$$

These sets define how transmitters interact based on their relative distances and orientations, essential for modeling the RTS/CTS mechanism's impact on network interference. Consider a transceiver pair with its transmitter at the origin  $o$  and receiver at  $(d, 0)$ , and another pair with its transmitter at  $(r \cos \beta, r \sin \beta)$  and receiver oriented at  $\theta$ . The set  $S_1$  includes  $(r, \beta, \theta)$  where the second pair's transmitter is within the physical carrier sensing region of the first pair.  $S_2$  includes  $(r, \beta, \theta)$  where the second pair's transmitter is within the RTS/CTS region influenced by the first pair's receiver.  $S_3$  covers  $(r, \beta, \theta)$  where the first pair's transmitter is within the RTS/CTS region created by the second pair's transmitter. The union  $S_1 \cup S_2 \cup S_3$  indicates the conditions under which at least one transceiver pair will be suppressed from transmission.

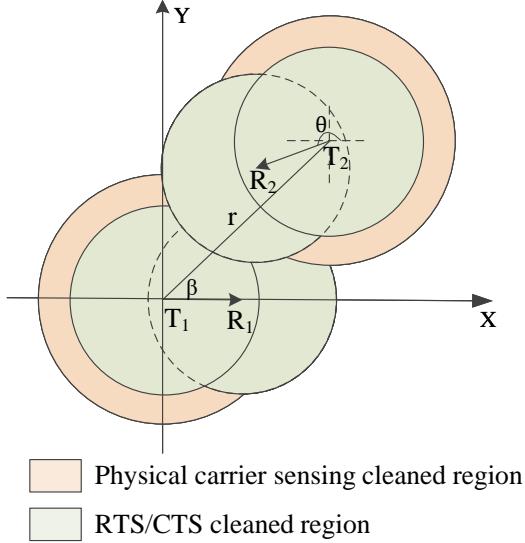


Fig. 3. Spatial distribution of two transceiver pairs showing their respective exclusion regions.

- Combined sensing area  $V(r, \beta, \theta)$ : Describes the combined area of the physical and virtual carrier sensing regions affected by two transceiver pairs separated by distance  $r$ , with relative phase difference  $\beta$ , and orientations  $0$  and  $\theta$ , respectively (see Figure 3).

These formulations enable us to rigorously calculate the density of active transmitters and the resultant mean interference, accounting for both the geometrical layout and the protective measures instituted by the RTS/CTS mechanism.

#### A. Type I Dual-Zone Hard-Core Process

1) *Node Intensity*: In analyzing the node intensity and the mean interference for the Type I dual-zone hard-core process, we use the concept of Palm measure. This measure is particularly useful for describing the statistical properties of a point process observed from a given point. For a stationary point process  $\Psi$  with a finite, non-zero intensity  $\lambda$ , and defined on a measurable space  $(\mathbb{M}, \mathcal{M})$  with  $\mathbb{M}$  representing all possible simple point patterns and  $\mathcal{M}$  a  $\sigma$ -algebra over  $\mathbb{M}$ , the Palm distribution can be expressed as [38, Sec. 8.2.1]

$$P_o(Y) = \frac{1}{\lambda v_d(B)} \int_{\mathbb{M}} \sum_{x \in \psi \cap B} \mathbb{1}_Y(\psi_{-x}) P(d\psi), \quad Y \in \mathcal{M}. \quad (10)$$

Here,  $v_d$  is the Lebesgue measure,  $B$  an arbitrary Borel set with positive measure, and  $\psi_{-x} = \{y \in \psi : y - x\}$  denotes the translation of  $\psi$  by  $-x$ .

Focusing on the Type I dual-zone hard-core process, where  $\Phi$  is a dependently thinned version of the original PPP  $\Phi_{\text{PPP}}$ , we consider the Palm measure  $P_o^{(\Phi_{\text{PPP}})}$  of  $\Phi_{\text{PPP}}$ . The node intensity for this process is

$$\lambda = \lambda_p P_o^{(\Phi_{\text{PPP}})}(o \in \Phi) = \lambda_p e^{-\lambda_p V_o}, \quad (11)$$

where  $V_o$  is the area of the RTS/CTS exclusion region, factoring in both physical and virtual carrier sensing ranges.

This expression suggests that the node intensity  $\lambda$  of the thinned process equals the intensity  $\lambda_p$  of the original PPP adjusted by the void probability within the exclusion region defined by  $V_o$ . This probability highlights how likely a space is free from other nodes, thereby not preventing the considered node's transmission. When  $\lambda_p$  is relatively low,  $\lambda$  increases linearly with  $\lambda_p$ . Conversely, as  $\lambda_p$  increases beyond a certain threshold,  $\lambda$  begins to decrease, underscoring the overly conservative nature of the Type I thinning in dense networks. The maximum efficiency, where  $\lambda$  peaks, is achieved when  $\lambda_p$  is optimally balanced against the exclusion region's area, which occurs at  $\lambda_p = 1/V_o$ .

2) *Mean Interference*: This subsection derives the mean interference in the Type I dual-zone hard-core process. Without loss of generality, we consider a scenario where the transmitter is located at the origin  $o$  and the receiver is at  $z_o = (d \cos \theta_o, d \sin \theta_o)$ . We specifically condition our analysis on the event  $(o, \theta_o) \in \tilde{\Phi}$ , ensuring that the point  $(o, \theta_o)$ , representing the signal source, does not contribute to the interference calculations. This exclusion is critical as the signal from the considered transceiver pair does not constitute interference to itself.

We employ the *reduced* Palm measure, denoted as  $P_{(o, \theta_o)}^!$ , associated with the marked point process  $\tilde{\Phi}$  for our interference analysis. The expectation under this measure, denoted by  $E_{(o, \theta_o)}^!$ , helps calculate the expected interference  $I_{z_o}$ . Notably, the mean interference calculation is independent of  $\theta_o$ , allowing us to simplify our notation to  $E_{(o, 0)}^!(I_{z_o})$ . Conditioning on  $\theta_o$  provides necessary flexibility for deriving broader results applicable in varying geometrical configurations.

**Theorem 1.** *The mean interference  $I_{z_o}$  that a receiver experiences in the Type I dual-zone hard-core process is*

$$E_{(o, 0)}^!(I_{z_o}) = \frac{\lambda_p^2 P_t}{2\pi\lambda} \int_0^\infty \int_0^{2\pi} \int_0^{2\pi} l(\sqrt{r^2 - 2rd \cos \beta + d^2}) k(r, \beta, 0, \theta) r d\theta d\beta dr, \quad (12)$$

where the kernel function  $k(r, \beta, 0, \theta)$  is defined as:

$$k(r, \beta, 0, \theta) = \begin{cases} 0 & \text{if } (r, \beta, \theta) \in S_1 \cup S_2 \cup S_3 \\ \exp(-\lambda_p V(r, \beta, \theta)) & \text{otherwise.} \end{cases} \quad (13)$$

*Proof.* See Appendix A.  $\square$

This theorem encapsulates the core analytical challenge of calculating the interference by integrating over the spatial variables and taking into account the non-uniform distribution of potential interfering transmitters induced by the dual-zone exclusion mechanism. The sets  $S_1$ ,  $S_2$ , and  $S_3$  delineate regions where interfering transmitters are either within the physical sensing or the RTS/CTS cleaned regions, thus significantly affecting the resultant interference patterns. By incorporating these spatial dependencies, we offer a comprehensive model that accurately reflects the complexities introduced by the RTS/CTS mechanism in managing interference.

#### B. Type II Dual-Zone Hard-Core Process

1) *Node Intensity*: In the Type II dual-zone hard-core process, each transceiver pair is marked with a time stamp

$t$ , which significantly influences the retention of transmitters within the process. For a given transmitter at a specific location, those transmitters within its induced guard region and having a time stamp greater than  $t$  are removed, simulating a priority system based on the time of initiation.

This mechanism leads to a situation where the retention probability of a transceiver pair is dependent on its time stamp in relation to others within its vicinity. Specifically, the probability that a given transceiver pair with time stamp  $t$  is retained is  $e^{-\lambda_p t V_o}$ , reflecting the exponential thinning of the process based on the area of the guard region  $V_o$  and the density of transmitters  $\lambda_p$ .

The overall intensity  $\lambda$  of the Type II process is then derived by integrating this retention probability over all possible values of  $t$  from 0 to 1, representing a uniform distribution of time stamps:

$$\begin{aligned}\lambda &= \lambda_p P_o^{(\Phi_{\text{PPP}})} (o \in \Phi) \\ &= \lambda_p \int_0^1 P((o, \theta_o) \in \tilde{\Phi} | (o, \theta_o, m_o) \in \tilde{\Phi}_o, m_o = t) dt \\ &= \lambda_p \int_0^1 e^{-\lambda_p t V_o} dt \\ &= \frac{1}{V_o} (1 - e^{-\lambda_p V_o}),\end{aligned}\quad (14)$$

where the integral simplifies to the expression seen in the final line. Here,  $V_o$  is defined as before in (9) and represents the combined area of both the RTS and CTS regions influencing the potential for a transmitter to be active.

This resulting intensity  $\lambda$  provides insight into how the system behaves as  $\lambda_p$  increases. Notably,  $\lambda$  increases monotonically with  $\lambda_p$  and asymptotically approaches the limit  $\frac{1}{V_o}$  as  $\lambda_p \rightarrow \infty$ . This limit signifies that the maximum possible density of active transmitters is inversely proportional to the area of the exclusion region, regardless of the initial density  $\lambda_p$  of potential transmitters. Thus, in highly dense environments, the process effectively becomes limited by the spatial constraints imposed by  $V_o$ , rather than the initial density of transmitters.

2) *Mean Interference:* Building upon the framework used in the Type I process, we derive the mean interference experienced by the typical receiver in the Type II dual-zone hard-core process. This process accounts for time stamps that influence transmitter retention, providing a dynamic approach to managing interference.

**Theorem 2.** *The mean interference  $I_{z_o}$  at the typical receiver located at  $z_o$  in the Type II process is given by*

$$E_{(o,0)}^!(I_{z_o}) \frac{\lambda_p^2 P_t}{2\pi\lambda} \int_0^\infty \int_0^{2\pi} \int_0^{2\pi} l(\sqrt{r^2 - 2rd\cos\beta + d^2}) k(r, \beta, 0, \theta) r d\theta d\beta dr, \quad (15)$$

where

$$k(r, \beta, 0, \theta) = \begin{cases} 0 & (r, \beta, \theta) \in S_1 \cup (S_2 \cap S_3) \\ 2\eta(V) & (r, \beta, \theta) \in S_1 \cap S_2 \cap S_3 \\ \eta(V) & \text{otherwise,} \end{cases} \quad (16)$$

in which  $V = V(r, \beta, \theta)$ , and

$$\eta(V) = \frac{V_o e^{-\lambda_p V} - V e^{-\lambda_p V_o} + V - V_o}{\lambda_p^2 (V - V_o) V V_o}. \quad (17)$$

*Proof.* See Appendix B.  $\square$

In the theorem, the function  $\eta(V)$  accounts for the probability modifications due to the presence of other potential transmitters within these regions, considering both the original and the exclusion region effects. The complexity in  $k(r, \beta, 0, \theta)$  arises from the necessity to consider overlapping guard regions, where multiple transmitters might influence the retention of a given pair. The conditional structure  $S_1 \cap S_2 \cap S_3$  indicates configurations where neither transmitter lies within the immediate guard region of the other, requiring an adjusted probability function  $2\eta(V)$  to account for the doubled effect when both are likely to be retained. The otherwise clause captures all other configurations, applying a single  $\eta(V)$  adjustment.

For the case where  $R_{\text{tx}} + d < R_{\text{cs}}$ , meaning the RTS/CTS cleaned region is contained within the physical carrier sensing cleaned region, we have  $S_2 \subset S_1$  and  $S_3 \subset S_1$ . In this scenario, the dual-zone hard-core process simplifies to the Matérn hard-core process. Consequently, the function  $k(r, \beta, 0, \theta)$  aligns with the results presented in [13] for both Type I and Type II Matérn hard-core processes.

#### IV. SUCCESS PROBABILITY

Calculating the success probability  $p_s(\gamma_o)$  of transmissions within spatial point process models requires the probability generating functional (PGFL). However, it is most likely impossible to find an analytical expression for the PGFL of the complex dual-zone hard-core process, which prevents an exact calculation of the success probability.

To address this limitation, prior research has introduced an alternative approach that uses an asymptotic gain to approximate the success probability in cellular networks relative to that of a PPP [15]–[18]. It was demonstrated that the proposed approximation approach is effective for cellular models, where the intended transmitter is the closest point of a point process. However, its application to bipolar models has not been explored. The interference characteristics in bipolar models, especially under a hard-core constraint, differ substantially due to their more structured point process. In this work, we extend the applicability of this approximation method for the first time to the bipolar dual-zone hard-core model, where the distance to the associated transmitter is fixed.

The proposed method involves first determining the asymptotic gain  $G$  for the point process model  $\phi$ . The success probability for  $\phi$ , denoted by  $P_\phi(T)$ , is then approximated using the success probability of a PPP model with the nearest transmitter association,  $P_{\text{PPP}}(T)$ , adjusted by  $G$ :

$$P_\phi(T) \approx P_{\text{PPP}} \left( \frac{T}{G} \right). \quad (18)$$

The homogeneous independent Poisson (HIP) model with the nearest transmitter association is used as a reference for deriving  $P_{\text{PPP}}(T)$  due to its analytical tractability and

simplicity. The HIP model assumes no intra-tier or inter-tier dependencies, and its signal-to-interference ratio (SIR) distribution is independent of node density or transmit power levels. This invariance allows the use of a single-tier framework to compute the Mean Interference-to-Signal Ratio (MISR), which is central to deriving the asymptotic gain.

The asymptotic gain  $G$ , as defined in [15], is a measure of the relative SIR characteristics of a PPP and the point process model  $\phi$ . It is expressed as the ratio of the MISR of a PPP to that of  $\phi$ :

$$G = \frac{\text{MISR}_{\text{PPP}}}{\text{MISR}_{\phi}}. \quad (19)$$

Here,  $\text{MISR}_{\phi}$  quantifies the average interference relative to the received signal strength in the point process model  $\phi$ . For the HIP model with the nearest transmitter association, the MISR is [15]

$$\text{MISR}_{\text{PPP}} = \frac{2}{\alpha - 2}. \quad (20)$$

The ratio  $G$  reflects the variation in system performance under two different interference conditions. In particular, when  $G$  is large, it indicates that the interference in the first condition has a more significant impact relative to the signal, potentially leading to a degradation in system performance.

We will next provide a detailed exposition on deriving the asymptotic gain  $G$  for both Type I and Type II dual-zone hard-core processes.

**Lemma 1.** *The asymptotic gain for Type I and Type II of the dual-zone hard-core processes is given by*

$$G_{\text{DZHCP}} = \frac{4\pi\lambda A}{(\alpha - 2)\lambda_p^2 r_0^\alpha} \left( \int_0^\infty \int_0^{2\pi} \int_0^{2\pi} l(\sqrt{r^2 - 2rd\cos\beta + d^2}) k(r, \beta, 0, \theta) r d\theta d\beta dr \right)^{-1}, \quad (21)$$

where  $\lambda_p$  is the intensity of the underlying PPP,  $k(r, \beta, 0, \theta)$  is given by (13) for Type I of the dual-zone hard-core process and given by (16) for Type II of the dual-zone hard-core process.

*Proof.* Calculating the MISR for the Type I and Type II dual-zone hard-core processes is considerably more complex than a PPP model. Unlike the PPP, these models do not benefit from the equivalence properties provided by Slivnyak's theorem, which simplifies analysis by asserting that conditioning on having a point at  $o$  is equivalent to simply adding  $o$  to the PPP. For the dual-zone hard-core processes, however, we must specifically evaluate the mean interference at a receiver's location, denoted by  $\mathbb{E}_{(o,0)}^! [I_{z_o}]$ , conditioned on having a transceiver pair with the receiver positioned at that location.

For the Type I and Type II processes, the mean interference at the receiver can be calculated using Theorem 1 and Theorem 2, respectively. The MISR for these models is then determined by dividing the mean interference by the average received power at the reference distance  $r_0$ , which is

$$\text{MISR}_{\text{DZHCP}} = \frac{\lambda_p^2 r_0^\alpha}{2\pi\lambda A} \int_0^\infty \int_0^{2\pi} \int_0^{2\pi} l(\sqrt{r^2 - 2rd\cos\beta + d^2}) k(r, \beta, 0, \theta) r d\theta d\beta dr, \quad (22)$$

where  $k(r, \beta, 0, \theta)$  is defined by (13) for Type I and by (16) for Type II of the dual-zone hard-core process.

Applying the definition of asymptotic gain from (19), we can then derive the corresponding results as outlined in Lemma.  $\square$

Considering a Rayleigh fading scenario where the expected value of the channel fading coefficient  $h$  is  $\mathbb{E}[h] = 1$ , the success probability for the HIP model with nearest transmitter association for a given SIR threshold  $T$  is [11]

$$P_{\text{PPP}}(T) = \left( 1 + T^{\frac{2}{\alpha}} \int_{T^{-\frac{2}{\alpha}}}^\infty \frac{1}{1+t^{\frac{\alpha}{2}}} dt \right)^{-1}. \quad (23)$$

This expression encapsulates the cumulative distribution function (ccdf) of the SIR. For a path-loss exponent  $\alpha = 4$ , the equation simplifies to

$$P_{\text{PPP}}(T) = \frac{1}{1 + \sqrt{T} \arctan(\sqrt{T})}. \quad (24)$$

For any point process model  $\phi$  characterized by an asymptotic gain  $G$ , as derived in Lemma 1, the success probability can be approximated by scaling the threshold  $T$  by  $G$ :

$$P_{\phi}(T) \approx P_{\text{PPP}} \left( \frac{T}{G} \right) = \left( 1 + \left( \frac{T}{G} \right)^{\frac{2}{\alpha}} \int_{(\frac{T}{G})^{-\frac{2}{\alpha}}}^\infty \frac{1}{1+t^{\frac{\alpha}{2}}} dt \right)^{-1}. \quad (25)$$

Specifically, for  $\alpha = 4$ , the success probability for model  $\phi$  is approximated by

$$P_{\phi}(T) \approx \left( 1 + \sqrt{\frac{T}{G}} \arctan \left( \sqrt{\frac{T}{G}} \right) \right)^{-1}. \quad (26)$$

The asymptotic gain  $G$  varies depending on the spatial characteristics and interference patterns specific to the point process model under consideration. By selecting the appropriate asymptotic gain for each model, interference analysis can be precisely tailored to the unique spatial dynamics of each process. The provided approach enables a rigorous estimation of the success probability for Type I and Type II dual-zone hard-core processes, offering practical insights into their performance relative to standard PPP models. For Type I and Type II of the dual-zone hard-core processes, we get the following theorem.

**Theorem 3.** *The success probability for a receiver in the Type I and Type II dual-zone hard-core processes is given by*

$$P_{\text{DZHCP}}(T) = \left( 1 + \left( \frac{T}{G_{\text{DZHCP}}} \right)^{\frac{2}{\alpha}} \int_{(\frac{T}{G_{\text{DZHCP}}})^{-\frac{2}{\alpha}}}^\infty \frac{1}{1+t^{\frac{\alpha}{2}}} dt \right)^{-1}, \quad (27)$$

where  $G_{\text{DZHCP}}$  is the asymptotic gain calculated as outlined in Lemma 1 and defined in (21).

For a path-loss exponent  $\alpha = 4$ , a more specific form of the approximated success probability can be found for the same processes:

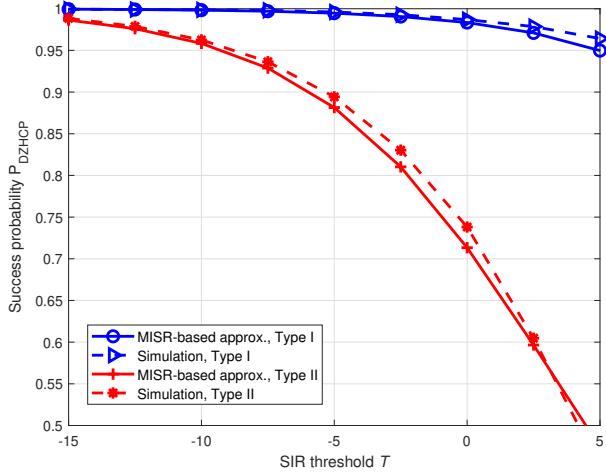


Fig. 4. SIR ccdf for the dual-zone hard-core process and MISR-based approximation for  $\lambda_p = 1 \times 10^{-4} \text{m}^{-2}$ .

**Corollary 1.** *The success probability for a receiver in the Type I and Type II dual-zone hard-core processes, with  $\alpha = 4$ , simplifies to*

$$P_{\text{DZHCP}}(T) \approx \left( 1 + \sqrt{\frac{T}{G_{\text{DZHCP}}}} \arctan \left( \sqrt{\frac{T}{G_{\text{DZHCP}}}} \right) \right)^{-1}, \quad (28)$$

where  $G_{\text{DZHCP}}$  is defined in (21).

These formulas model the success probability by adapting the classical success probability formula for a PPP, modified by the asymptotic gain which accounts for the reduced interference due to the dual-zone hard-core process. The success probability expression provides a quantifiable measure of how these strategies improve network performance relative to a standard PPP, particularly under various environmental and operational conditions dictated by the parameter  $\alpha$ . The average interference in the network is directly proportional to the transmission power, as established in Theorems 1 and 2. However, since both the signal power and interference power scale proportionally with transmission power, the SIR remains unchanged. As a result, the success probability is not affected by variations in transmission power, as confirmed by Theorem 3.

To assess the accuracy of the MISR approximation for the bipolar case, we compare its results with simulation data in this work. Figures 4 and 5 present the success probability under varying densities and for commonly used values of  $T$ , showing that the proposed MISR approximation yields accurate results for all practical success probabilities. These results validate the applicability of the MISR approximation for the bipolar dual hard-core model, despite the differences in interference characteristics when compared to cellular networks.

## V. NUMERICAL RESULTS

The numerical results are obtained from the analytical results we have derived. The default configurations of system model are as follows (also see Table I). Unless otherwise

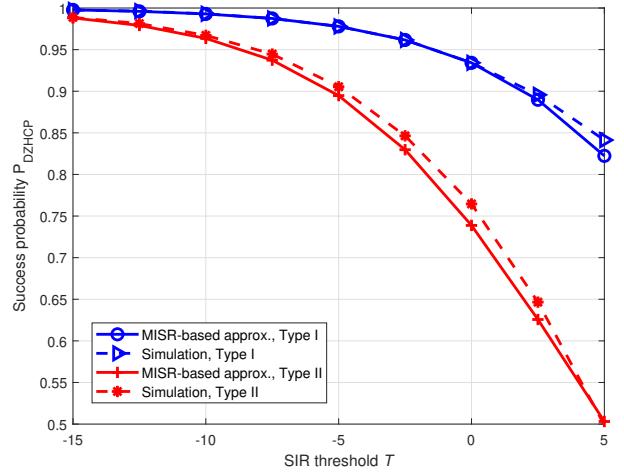


Fig. 5. SIR ccdf for the dual-zone hard-core process and MISR-based approximation for  $\lambda_p = 5 \times 10^{-5} \text{m}^{-2}$ .

TABLE I  
SYSTEM PARAMETERS

Symbol	Description	Value
$\lambda_p$	Intensity of potential transmitters	$1 \times 10^{-5} \text{m}^{-2}$
$\lambda$	Intensity of the Type I or II process	N/A
$R_{\text{tx}}$	Virtual carrier sensing radius	100m
$R_{\text{cs}}$	Physical carrier sensing radius	$1.2R_{\text{tx}}$
$d$	Distance between transmitter and receiver	$0.8R_{\text{tx}}$
$P_t$	Transmit power	20dBm
$T$	SIR threshold	0 dB
$\alpha$	Path loss exponent	3.5

specified, the intensity of potential transmitters is set as  $\lambda_p = 1 \times 10^{-5} \text{m}^{-2}$  and the virtual carrier sensing radius is  $R_{\text{tx}} = 100\text{m}$ . The physical carrier sensing radius is always set as  $R_{\text{cs}} = 1.2R_{\text{tx}}$  and the distance between transmitter and receiver is set as  $d = 0.8R_{\text{tx}}$ . The transmit power is set as  $P_t = 20\text{dBm}$  (0.1W), and the path loss model is  $l(r) = Ar^{-\alpha}$  with  $\alpha = 3.5$  and  $A = 0.01$ . In the following figures, the legend entries ‘Type I’ and ‘Type II’ represent the RTS/CTS handshake protocol as modeled by the Type I and Type II dual-zone hard-core processes, respectively. Similarly, the entries ‘CSMA I’ and ‘CSMA II’ correspond to the CSMA protocol, modeled using the Type I and Type II Matérn hard-core processes, respectively. We analyze a key performance metric, defined as the product of the success probability and the active transmitter density. This metric quantifies the spatial density of successful transmissions within a given discrete time interval, effectively capturing both the likelihood of successful transmission and the intensity of transmission attempts. By incorporating both spatial and temporal dimensions, it provides a measure of network throughput.

Figure 6 illustrates the relationship between the active node density  $\lambda$  and the original PPP intensity  $\lambda_p$  under different configurations of the RTS/CTS mechanism, with fixed virtual carrier sensing range  $R_{\text{tx}}$  and physical carrier sensing range  $R_{\text{cs}}$ . For both CSMA I and Type I models, we observe an initial increase in active node density  $\lambda$  as the original PPP intensity

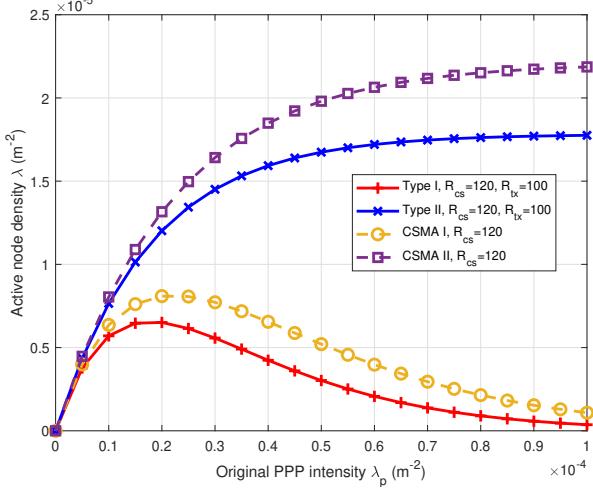


Fig. 6. Active node density  $\lambda$  as a function of the original PPP intensity  $\lambda_p$  for various RTS/CTS configurations.

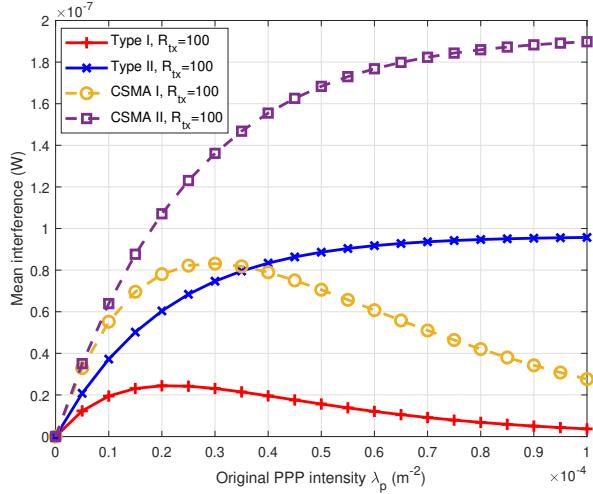


Fig. 7. Mean interference as a function of the original PPP intensity  $\lambda_p$  under different schemes. In the Type II model, protection regions are established around both transmitters and receivers, reducing the likelihood of interferers being located near the receiver. This additional protection results in lower mean interference compared to CSMA I, even as network node density increases.

$\lambda_p$  increases, followed by a subsequent decrease. This trend reaches its peak when  $\lambda_p = \frac{1}{V_o}$ , where  $V_o$  represents the area of the exclusion region, yielding a maximum value of  $\frac{1}{V_o} e^{-1}$ . This behavior reflects the impact of the protection region, where the density of active nodes is initially enhanced by a higher node availability but later constrained as interference effects intensify. In contrast, for the CSMA II and Type II models, active node density  $\lambda$  increases continuously with the original PPP intensity  $\lambda_p$ , eventually approaching the reciprocal of the exclusion region. This gradual increase indicates that these models allow for higher active node density even at larger values of  $\lambda_p$ , suggesting a more effective management of interference as  $\lambda_p$  grows.

Figure 7 illustrates the variation in mean interference as the original PPP intensity  $\lambda_p$  changes, with a fixed virtual carrier

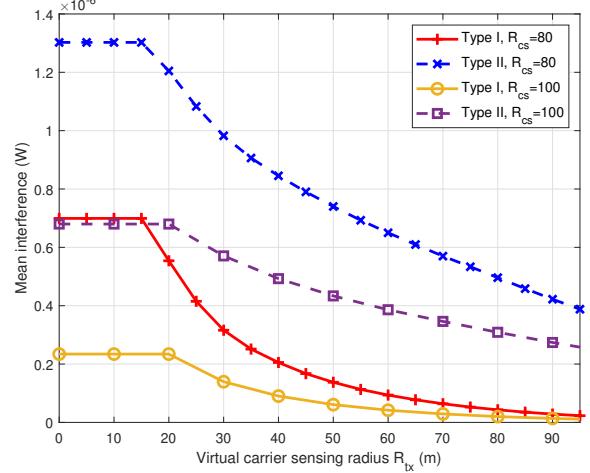


Fig. 8. Mean interference as a function of virtual carrier sensing radius  $R_{tx}$  with fixed physical carrier sensing radius  $R_{cs}$  and distance  $d$ .

sensing range  $R_{tx}$ , physical carrier sensing radius  $R_{cs}$ , and transceiver distance  $d$ . For CSMA I and Type I schemes, the mean interference initially increases with  $\lambda_p$  before eventually decreasing, mirroring the behavior observed in their active node densities, with both metrics peaking at the same  $\lambda_p$  value. A similar trend is observed for CSMA II and Type II schemes, where the mean interference follows the variation in active node density. It is observed that when  $\lambda_p < 3 \times 10^{-5} \text{ m}^{-2}$ , CSMA I exhibits a lower active node density than Type II, yet its mean interference remains higher. This seemingly counterintuitive result stems from the fundamental differences in the protection regions of the two mechanisms. In CSMA I, the absence of a receiver-side protection region allows interferers to be located closer to the receiver, increasing the probability of high interference. In contrast, the Type II model establishes protection zones around both the transmitter and receiver, effectively reducing the likelihood of nearby interferers. As a result, despite the lower active node density, CSMA I experiences a higher mean interference in this regime.

Figure 8 illustrates the relationship between mean interference and the virtual carrier sensing range  $R_{tx}$  under varying schemes with a constant physical carrier sensing radius  $R_{cs}$  and distance  $d$ . When  $R_{tx}$  is relatively small, the mean interference does not significantly change with increasing  $R_{tx}$ . This is because, at smaller values of  $R_{tx}$ , the protection region of the transmitter fully covers the protection region of the receiver, leading to an interference level that decreases as  $R_{cs}$  increases. However, as  $R_{tx}$  continues to increase and the transmitter's protection region no longer completely encompasses the receiver's protection region, the mean interference starts to decrease with increasing  $R_{tx}$ .

Figure 9 examines the relationship between mean interference and virtual carrier sensing range  $R_{tx}$  in Type II and CSMA II models, with a virtual carrier sensing radius  $R_{cs} = 1.2R_{tx}$ . In both models, as  $R_{tx}$  increases, indicating an expansion in protection region, the three CSMA II curves begin to converge, showing a similar trend. A similar convergence phenomenon is observed for the three curves in the Type

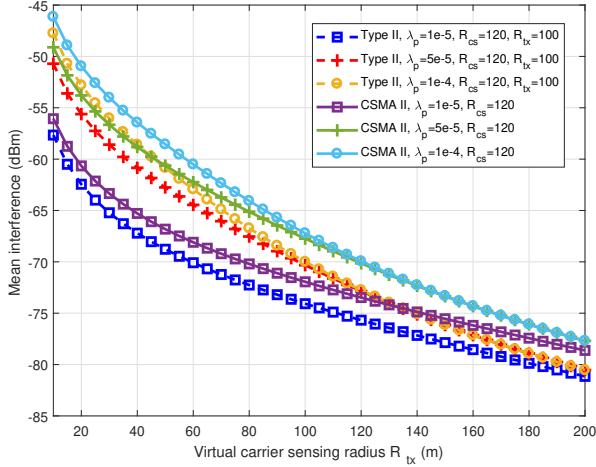


Fig. 9. Mean interference as a function of virtual carrier sensing radius  $R_{tx}$  for Type II and CSMA II models with physical carrier sensing radius  $R_{cs} = 1.2R_{tx}$ .

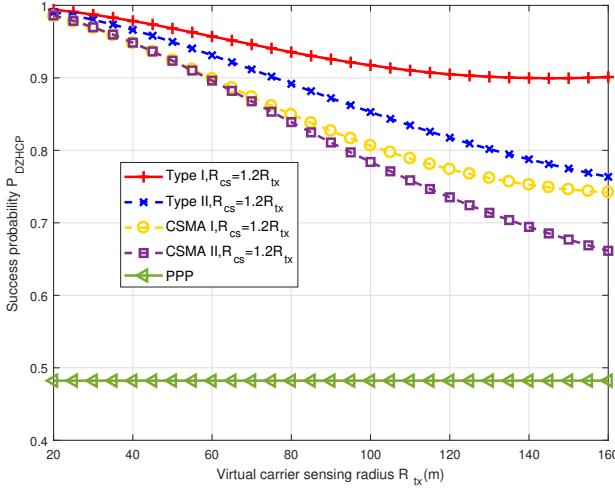


Fig. 10. Success probability as a function of virtual carrier sensing radius  $R_{tx}$  with fixed ratio  $d : R_{tx} : R_{cs} = 0.8 : 1 : 1.2$ .

II model. This pattern suggests that for the Type II hard-core point processes, as the protection region expands, the mean interference level becomes more dependent on the size of the protection region rather than the original PPP intensity.

Figure 10 illustrates the relationship between success probability  $P_{DZHP}$  and the virtual carrier sensing range  $R_{tx}$ , under the fixed ratio  $d : R_{tx} : R_{cs} = 0.8 : 1 : 1.2$ . As  $R_{tx}$  increases, there is a noticeable decrease in the success probability  $P_{DZHP}$ . This trend can be explained by two competing effects of increasing  $R_{tx}$ : while a larger  $R_{tx}$  helps to reduce mean interference by expanding the protection region, it simultaneously weakens the received signal strength, which negatively impacts the success probability. This phenomenon highlights that, under proportional scaling, the rate of reduction in mean interference is slower than the rate of decrease in received signal strength, leading to an overall reduction in success probability as  $R_{tx}$  grows.

Figure 11 shows the relationship between the success prob-

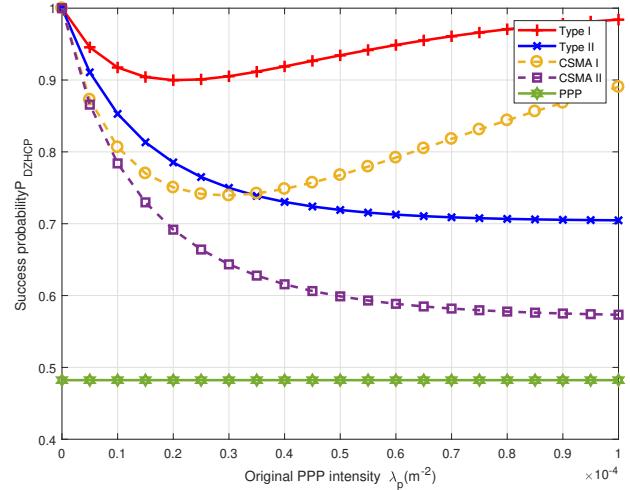


Fig. 11. Success probability  $P_{DZHP}$  as a function of PPP intensity  $\lambda_p$ .

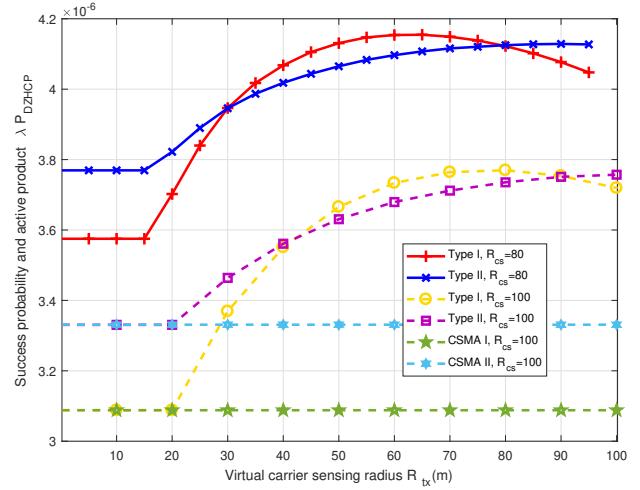


Fig. 12. Success probability and active density product as a function of  $R_{tx}$  (fixed  $R_{cs}$  and  $d$ ).

ability  $P_{DZHP}$  and the original PPP intensity  $\lambda_p$  for different network models. For the Type I and CSMA I models, the success probability  $P_{DZHP}$  initially decreases with an increase in  $\lambda_p$ , then increases after reaching a certain point. This behavior can be explained by the effect of mean interference, which first rises with increasing  $\lambda_p$  but then diminishes as the density becomes very high. On the other hand, for the Type II and CSMA II models, the success probability  $P_{DZHP}$  generally decreases as  $\lambda_p$  increases, but this decrease is not linear. Specifically, at lower values of  $\lambda_p$ , the success probability may drop quickly, while at higher values, the rate of decrease slows down, showing a gradual decline. This pattern suggests that for these models, the success probability asymptotically stabilizes as the PPP intensity increases.

Figure 12 illustrates the relationship between the product of success probability and active density with respect to the virtual carrier sensing range  $R_{tx}$ , under fixed values of  $R_{cs}$  and  $d$ . When  $R_{tx}$  is small, this metric remains relatively constant as  $R_{tx}$  increases. This stability occurs because, at lower  $R_{tx}$  values, the protection region of the transmitter

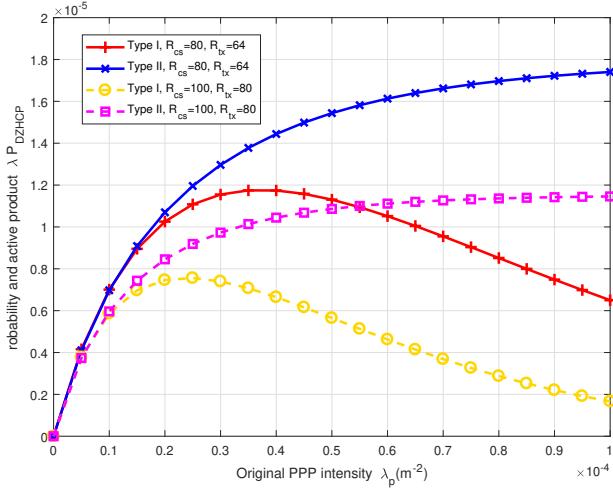


Fig. 13. Success probability and active density product as a function of the original PPP intensity  $\lambda_p$ .

completely covers the receiver's protection region, leading to no interference fluctuation. However, as  $R_{tx}$  continues to increase, there comes a point where the transmitter's protection region no longer fully covers the receiver's region. For the Type I model, beyond this point, the metric first increases with  $R_{tx}$  and then decreases. Notably, the maximum value for both cases ( $R_{cs} = 80m$  and  $R_{cs} = 100m$ ) occurs around  $R_{tx} = 0.8R_{cs}$ , indicating that the network throughput is optimal at this proportion for this type. On the other hand, for the Type II model, as  $R_{tx}$  continues to increase, the metric would monotonically increase and the rate of this increase diminishes. This observation indicates that for the Type II model, the ability of  $R_{tx}$  to effectively enhance network throughput diminishes as it increases. By comparing the performance of CSMA with that of the RTS/CTS mechanism, it becomes evident that the RTS/CTS mechanism not only mitigates interference but also significantly enhances the overall network throughput.

Figure 13 illustrates the relationship between the product of success probability and active density with respect to the intensity of the original PPP  $\lambda_p$ . As shown in the figure, Type I networks exhibit a trend where this metric initially increases with the increase in  $\lambda_p$ , reaching a peak, and then decreases as  $\lambda_p$  continues to grow. This pattern suggests an optimal network density where the product of success probability and active density is maximized. For Type II, as  $\lambda_p$  continues to increase, the metric would monotonically decrease and then gradually stabilize, indicating that the system reaches a saturation point where further increases in  $\lambda_p$  have a diminishing impact on the system throughput.

## VI. CONCLUSION

In this paper, we developed marked point process models to analyze the spatial distribution of transceivers under the RTS/CTS handshake mechanism in WLANs. Our study quantifies the impact of carrier sensing ranges and node density on mean interference and transmission success probability, offering insights into network performance across different

spatial configurations. The results reveal the crucial role of transceiver placement and handshake mechanisms in managing interference. Our findings demonstrate the potential of correlated point process modeling in optimizing wireless network design.

Future research could focus on improving the Type II model to address its underestimation of active transmitter density, thereby enhancing its accuracy in evaluating interference and network performance. Another promising direction is to extend the model to incorporate directional transmission and reception empowered by beamforming, which play a crucial role in modern wireless networks but are not captured by the current omnidirectional transmission assumption.

## APPENDIX A PROOF OF THEOREM 1

We define  $\tilde{\mathbb{M}}$  as the set of all marked point processes. The interference experienced by the receiver located at  $z_o$  is calculated by summing the powers received from all other transmitters in the network, where each transmitter is specified by its location  $x$  and orientation  $\theta$  within the point process  $\tilde{\Phi}$ .

The expectation of interference,  $E_{(o, \theta_o)}^!(I_{z_o})$ , under the reduced Palm probability measure  $P_{(o, \theta_o)}^!$ , conditioned on a transmitter at the origin  $o$  with orientation  $\theta_o$ , is:

$$\begin{aligned} E_{(o, \theta_o)}^!(I_{z_o}) &= E_{(o, \theta_o)}^! \left( \sum_{(x, \theta) \in \tilde{\Phi}} P_t l(|x - z_o|) \right) \\ &= \lambda P_t \int_{\mathbb{R}^2 \times [0, 2\pi]} l(|x - z_o|) \mathcal{K}_{\theta_o}(d(x, \theta)). \end{aligned} \quad (29)$$

The reduced second-order factorial measure  $\mathcal{K}_{\theta_o}(B \times L)$  for regions  $B \subset \mathbb{R}^2$  and  $L \subset [0, 2\pi]$  is defined as

$$\mathcal{K}_{\theta_o}(B \times L) = \frac{1}{\lambda} \int_{\tilde{\mathbb{M}}} \tilde{\varphi}(B \times L) P_{(o, \theta_o)}^!(d\tilde{\varphi}), \quad (30)$$

where  $\tilde{\varphi}(B \times L)$  counts the points in  $\tilde{\Phi}$  that fall within  $B$  and  $L$ . This measure represents the expected number of points in  $B \times L$  under the Palm distribution  $P_{(o, \theta_o)}^!$ .

For the second-order factorial moment measure  $\alpha^{(2)}(B_1 \times B_2 \times L_1 \times L_2)$ , which represents pairs of points in specified spatial and mark regions, we have:

$$\begin{aligned} \alpha^{(2)}(B_1 \times B_2 \times L_1 \times L_2) &= \\ &\frac{\lambda^2}{2\pi} \int_{B_1 \times L_1} \mathcal{K}_{\theta}((B_2 - x) \times L_2) d(x, \theta). \end{aligned} \quad (31)$$

The second-order product density,  $\varrho^{(2)}(x_1, x_2, \theta_1, \theta_2)$ , describes the likelihood of finding pairs of points with specific relative positions and marks. For a stationary process, we have:

$$\begin{aligned} \alpha^{(2)}(B_1 \times B_2 \times L_1 \times L_2) &= \int_{B_1 \times L_1} \\ &\left( \int_{(B_2 - x_1) \times L_2} \varrho^{(2)}(x_2, \theta_1, \theta_2) d(x_2, \theta_2) \right) d(x_1, \theta_1). \end{aligned} \quad (32)$$

This leads to:

$$\mathcal{K}_{\theta_o}(B \times L) = \frac{2\pi}{\lambda^2} \int_{B \times L} \varrho^{(2)}(x, \theta_o, \theta) d(x, \theta). \quad (33)$$

Thus, the mean interference  $E_{(o,\theta_o)}^!(I_{z_o})$  becomes

$$E_{(o,\theta_o)}^!(I_{z_o}) = \frac{2\pi P_t}{\lambda} \int_{\mathbb{R}^2 \times [0,2\pi]} l(|x - z_o|) \varrho^{(2)}(x, \theta_o, \theta) d(x, \theta). \quad (34)$$

This expression incorporates the dependencies in location and orientation, providing insight into interference distribution.

To derive the second-order product density  $\varrho^{(2)}(x, \theta_o, \theta)$ , we utilize its relation to the two-point Palm probability  $k(x, \theta_o, \theta)$ . This probability quantifies the likelihood that two transmitters, separated by a vector  $x$  and with orientation marks  $\theta_o$  and  $\theta$ , are both active (i.e., not thinned out by the RTS/CTS mechanism). The relationship is given by the following equation derived from stochastic geometry principles:

$$\varrho^{(2)}(x, \theta_o, \theta) = \frac{\lambda_p^2}{4\pi^2} k(x, \theta_o, \theta). \quad (35)$$

Here,  $k(x, \theta_o, \theta)$  acts as a modulation factor adjusting this density based on the spatial and orientational configuration of transmitters.

We consider a typical scenario where the orientation of the initial transmitter is normalized to zero,  $\theta_o = 0$ , simplifying the analysis without loss of generality. The typical receiver is hence positioned at  $z_o = (d, 0)$ . The interference experienced by this receiver,  $I_{z_o}$ , is then calculated as:

$$E_{(o,0)}^!(I_{z_o}) = \frac{\lambda_p^2 P_t}{2\pi \lambda} \int_{\mathbb{R}^2 \times [0,2\pi]} l(|x - z_o|) k(x, 0, \theta) d(x, \theta), \quad (36)$$

where  $l(|x - z_o|)$  denotes the path-loss function, and  $k(x, 0, \theta)$  modifies the intensity of points contributing to the interference based on their spatial and mark relationships.

This integral can be expressed in polar coordinates as

$$E_{(o,0)}^!(I_{z_o}) = \frac{\lambda_p^2 P_t}{2\pi \lambda} \int_0^\infty \int_0^{2\pi} \int_0^{2\pi} l(\sqrt{r^2 - 2rd \cos \beta + d^2}) k(r, \beta, 0, \theta) r d\theta d\beta dr. \quad (37)$$

Here,  $r$  is the radial distance from the origin,  $\beta$  is the angle difference from the direction to the receiver, and  $\theta$  is the mark of the interfering transmitter.

The function  $k(r, \beta, 0, \theta)$ , critical for this calculation, reflects the conditional probability that two transmitters, separated by distance  $r$  with phase angle difference  $\beta$  and marks 0 and  $\theta$ , do not violate the RTS/CTS protocol rules. Specifically, it evaluates to zero if the triple  $(r, \beta, \theta)$  falls within any of the sets  $S_1$ ,  $S_2$ , or  $S_3$ , indicating immediate exclusion due to proximity within the guard regions. Otherwise, the transmitters are retained with a probability of  $\exp(-\lambda_p V(r, \beta, \theta))$ , reflecting the void probability of the exclusion region  $V(r, \beta, \theta)$  for the PPP.

## APPENDIX B PROOF OF THEOREM 2

The derivation of (15) follows the same line as that of the Type I dual-zone hard-core process in Theorem 1 except for the difference in the probability that both of the transceiver

pairs are retained, namely,  $k(r, \beta, 0, \theta)$ , which is the probability that two transceiver pairs are both retained, is different from the Type I dual-zone hard-core process, since the time stamp marks are taken into consideration. We should consider not only the geometric relationship (see Figure 3), but also the relationship between the time stamp marks of the two transceiver pairs. When  $(r, \beta, \theta) \in S_1 \cup (S_2 \cap S_3)$ , at least one of the considered two transceiver pairs has to be removed after comparing their time stamp marks and thus  $k(r, \beta, 0, \theta)$  is zero. When  $(r, \beta, \theta) \in \overline{S_1} \cap \overline{S_2} \cap \overline{S_3}$ , there is no direct comparison between the time stamp marks of the considered two transceiver pairs. Let their time stamp marks be  $t_1$  and  $t_2$  respectively. So by separately considering the cases of  $t_1 \geq t_2$  and  $t_1 < t_2$ , we get

$$\begin{aligned} k(r, \beta, 0, \theta) &= \int_0^1 e^{-\lambda_p t_1 V_o} \left[ \int_0^{t_1} e^{-\lambda_p t_2 (V - V_o)} dt_2 \right] dt_1 \\ &\quad + \int_0^1 e^{-\lambda_p t_2 V_o} \left[ \int_0^{t_2} e^{-\lambda_p t_1 (V - V_o)} dt_1 \right] dt_2 \\ &= 2 \int_0^1 e^{-\lambda_p t_1 V_o} \left[ \int_0^{t_1} e^{-\lambda_p t_2 (V - V_o)} dt_2 \right] dt_1, \end{aligned} \quad (38)$$

which can be evaluated as  $2\eta(V)$ . Here,  $\int_0^{t_1} e^{-\lambda_p t_2 (V - V_o)} dt_2$  is the conditional probability that, when the first transceiver pair is marked by  $t_1$ , there is no other transceiver pair of mark smaller than  $t_2 < t_1$  lying in the region covered by the second transceiver pair only;  $\int_0^{t_2} e^{-\lambda_p t_1 (V - V_o)} dt_1$  is analogously interpreted. When either  $(r, \beta, \theta) \in \overline{S_1} \cap S_2 \cap \overline{S_3}$  or  $(r, \beta, \theta) \in \overline{S_1} \cap \overline{S_2} \cap S_3$ , exactly one of the considered transmitter is within the RTS/CTS cleaned region of the other, and hence only one of the terms in (38) should be taken into account, resulting into  $k(r, \beta, 0, \theta) = \eta(V)$ .

## REFERENCES

- [1] Y. Zhong, W. Zhang, and M. Haenggi, "Stochastic analysis of the mean interference for the RTS/CTS mechanism," in *IEEE International Conference on Communications (ICC)*, 2014, pp. 1996–2001.
- [2] Y. Zhong and X. Ge, "Toward 6G Carbon-Neutral Cellular Networks," *IEEE Network*, vol. 38, no. 5, pp. 174–181, 2024.
- [3] IEEE 802.11 Working Group, "Wireless LAN medium access control (MAC) and physical layer (PHY) specifications," 1997.
- [4] A. Hasan and J. G. Andrews, "The guard zone in wireless ad hoc networks," *IEEE Transactions on Wireless Communications*, vol. 6, no. 3, pp. 897–906, 2007.
- [5] F. Tobagi and L. Kleinrock, "Packet switching in radio channels: part II—the hidden terminal problem in carrier sense multiple-access and the busy-tone solution," *IEEE Transactions on Communications*, vol. 23, no. 12, pp. 1417–1433, 1975.
- [6] S. Xu and T. Saadawi, "Does the IEEE 802.11 MAC protocol work well in multihop wireless ad hoc networks?" *IEEE Communications Magazine*, vol. 39, no. 6, pp. 130–137, 2001.
- [7] J. Weinmiller, H. Woesner, J. Ebert, and A. Wolisz, "Analyzing the RTS/CTS mechanism in the DFWMAC media access protocol for wireless LANs," *IFIP TC6 Workshop on Personal Wireless Communications*, 1995.
- [8] F. Cali, M. Conti, and E. Gregori, "IEEE 802.11 wireless LAN: capacity analysis and protocol enhancement," in *INFOCOM'98. Seventeenth Annual Joint Conference of the IEEE Computer and Communications Societies. Proceedings.*, vol. 1, 1998, pp. 142–149.
- [9] G. Bianchi, "Performance analysis of the IEEE 802.11 distributed coordination function," *IEEE Journal on Selected Areas in Communications*, vol. 18, no. 3, pp. 535–547, 2000.
- [10] K. Xu, M. Gerla, and S. Bae, "Effectiveness of RTS/CTS handshake in IEEE 802.11 based ad hoc networks," *Ad Hoc Networks*, vol. 1, no. 1, pp. 107–123, 2003.

[11] M. Haenggi, J. Andrews, F. Baccelli, O. Dousse, and M. Franceschetti, "Stochastic geometry and random graphs for the analysis and design of wireless networks," *IEEE Journal on Selected Areas in Communications*, vol. 27, no. 7, pp. 1029–1046, 2009.

[12] Y. Zhong, G. Mao, X. Ge, and F.-C. Zheng, "Spatio-temporal modeling for massive and sporadic access," *IEEE Journal on Selected Areas in Communications*, vol. 39, no. 3, pp. 638–651, 2020.

[13] M. Haenggi, "Mean Interference in Hard-Core Wireless Networks," *IEEE Communications Letters*, vol. 15, no. 8, pp. 792–794, 2011.

[14] B. Matérn, "Spatial variation," *Lecture Notes in Statistics*, vol. 36, 1986.

[15] M. Haenggi, "The Mean Interference-to-Signal Ratio and Its Key Role in Cellular and Amorphous Networks," *IEEE Wireless Communications Letters*, vol. 3, no. 6, pp. 597–600, 2014.

[16] R. K. Ganti and M. Haenggi, "Asymptotics and Approximation of the SIR Distribution in General Cellular Networks," *IEEE Transactions on Wireless Communications*, vol. 15, no. 3, pp. 2130–2143, 2016.

[17] S. S. Kalamkar and M. Haenggi, "Simple Approximations of the SIR Meta Distribution in General Cellular Networks," *IEEE Transactions on Communications*, vol. 67, no. 6, pp. 4393–4406, 2019.

[18] H. Wei, X. Xu, B. Lin, X. Lu, and P. Wang, "A Novel ASAPP Approach to Characterize the SIR Distribution in General Cellular Networks," *IEEE Wireless Communications Letters*, vol. 11, no. 7, pp. 1443–1447, 2022.

[19] Y. Yin, Y. Gao, S. Manzoor, and X. Hei, "Optimal RTS Threshold for IEEE 802.11 WLANs: Basic or RTS/CTS?" in *IEEE Smart-World, Ubiquitous Intelligence & Computing, Advanced & Trusted Computing, Scalable Computing & Communications, Cloud & Big Data Computing, Internet of People and Smart City Innovation (Smart-World/SCALCOM/UIC/ATC/CBDCom/IOP/SCI)*, 2019, pp. 1620–1625.

[20] T. T. Sari and G. Seçinti, "Chain RTS/CTS Scheme for Aerial Multihop Communications," in *IEEE 18th Annual Consumer Communications & Networking Conference (CCNC)*, 2021, pp. 1–2.

[21] M. Siew, S. Sharma, K. Guo, D. Cai, W. Wen, C. Joe-Wong, and T. Q. Quek, "Towards Effective Resource Procurement in MEC: A Resource Re-Selling Framework," *IEEE Transactions on Services Computing*, vol. 17, no. 1, pp. 82–97, 2024.

[22] F. Baccelli, J. Li, T. Richardson, S. Shakkottai, S. Subramanian, and X. Wu, "On optimizing CSMA for wide area ad hoc networks," *Queueing Systems*, vol. 72, no. 1-2, pp. 31–68, 2012.

[23] T. Tian, H. Iwai, and H. Sasaoka, "Statistical analysis of interference avoidance based on multi-frequency RTS/CTS cognitive radio," in *6th International ICST Conference on Cognitive Radio Oriented Wireless Networks and Communications (CROWNCOM)*, 2011, pp. 281–285.

[24] Y. Zhong, Z. Chen, T. Han, and X. Ge, "Spatio-Temporal Interference Correlation: Influence of Deployment Patterns and Traffic Dynamics," *IEEE Transactions on Communications*, 2024, accepted for publication.

[25] H. Nguyen, F. Baccelli, and D. Kofman, "A stochastic geometry analysis of dense IEEE 802.11 networks," in *IEEE 26th INFOCOM*, 2007.

[26] Z. Chen and Y. Zhong, "Characterizing stability regions in wireless networks with hard-core spatial constraints," *IEEE Wireless Communications Letters*, vol. 13, no. 8, pp. 2180–2184, 2024.

[27] K. V. Mishra, B. Shankar M.R., and B. Ottersten, "Stochastic-Geometry-Based Interference Modeling in Automotive Radars Using Matérn Hard-Core Process," in *IEEE Radar Conference (RadarConf20)*, 2020.

[28] Y. Zhu and S. Zhang, "Modeling and Analysis of Multi-UAV Networks Using Matérn Hard-Core Point Process," in *IEEE 23rd International Conference on High Performance Switching and Routing (HPSR)*, 2022, pp. 89–94.

[29] L. Kui, S. Huang, and Z. Feng, "Interference Analysis for Automotive Radar Using Matérn Hard-Core Point Process," in *IEEE/CIC International Conference on Communications in China (ICCC)*, 2021, pp. 817–822.

[30] G.-H. Li, D.-W. Yue, S.-N. Jin, and Q. Hu, "Multi-RIS Aided Multi-User mmWave MIMO System Modelled by Matérn Hard-Core Point Processes," *IEEE Transactions on Vehicular Technology*, vol. 73, no. 4, pp. 5988–5993, 2024.

[31] A. Busson, G. Chelius, and J.-M. Gorce, "Interference Modeling in CSMA Multi-Hop Wireless Networks," INRIA, Tech. Rep. RR-6624, 2009. [Online]. Available: <https://hal.inria.fr/inria-00316029v3>

[32] H. ElSawy and E. Hossain, "Modeling random CSMA wireless networks in general fading environments," in *IEEE International Conference on Communications (ICC)*, 2012, pp. 5457–5461.

[33] H. ElSawy and E. Hossain, "A Modified Hard Core Point Process for Analysis of Random CSMA Wireless Networks in General Fading Environments," *IEEE Transactions on Communications*, vol. 61, no. 4, pp. 1520–1534, 2013.

[34] A. S. Parihar, P. Swami, V. Bhatia, and Z. Ding, "Performance Analysis of SWIPT Enabled Cooperative-NOMA in Heterogeneous Networks Using Carrier Sensing," *IEEE Transactions on Vehicular Technology*, vol. 70, no. 10, pp. 10646–10656, 2021.

[35] G.-H. Shin, H.-K. Park, Y. Kim, S. Park, and T. Kwon, "Performance evaluation of LBT spectrum sharing via modified Matérn hardcore point process type III," in *International Conference on Information and Communication Technology Convergence (ICTC)*, 2017, pp. 492–494.

[36] F. Baccelli, P. Mühlthaler, and B. Blaszczyszyn, "Stochastic analysis of spatial and opportunistic ALOHA," *IEEE Journal on Selected Areas in Communications*, vol. 27, no. 7, pp. 1105–1119, 2009.

[37] Z. Tong and M. Haenggi, "Throughput analysis for full-duplex wireless networks with imperfect self-interference cancellation," *IEEE Transactions on Communications*, vol. 63, no. 11, pp. 4490–4500, 2015.

[38] M. Haenggi, *Stochastic Geometry for Wireless Networks*. Cambridge University Press, 2012.