

26 - pf of Thom's thm.

Note Title

5/13/2010

Thm: $\Omega_d \rightarrow \pi_d MO$ is an isomph.

(sketch) need an inv map.

given $\alpha \in \pi_d MO$

$$\alpha_N : S^{d+N} \longrightarrow MO(N)$$
$$BO(N)^{V_N}$$

$$BO(N) = \varprojlim_k \text{Gr}_N(\mathbb{R}^k)$$

$$\begin{array}{ccc}
 S^{d+N} & \longrightarrow & \text{Gr}_N(\mathbb{R}^\infty)^{\vee_N} \\
 \exists \quad \downarrow \alpha_N & & \uparrow \\
 & & \text{Gr}_N(\mathbb{R}^k)^{\vee_N} \\
 & \text{codim } N \uparrow \text{are rect} & \\
 & & \text{Gr}_N(\mathbb{R}^k)
 \end{array}$$

Then: $\alpha_N' \simeq \alpha_N''$

$$\begin{array}{c}
 \uparrow \\
 \alpha_N'' \notin \text{Gr}_N(\mathbb{R}^k)
 \end{array}$$

$$(\alpha_N'')^{-1}(\text{Gr}_N(\mathbb{R}^k)) = M \hookrightarrow S^{d+N}$$

$\underbrace{\phantom{\text{Gr}_N(\mathbb{R}^k)}}$
 $M = \text{smooth cpt}$
 $\text{codim } N$
 submfld
 embedded any
 $\text{from cpt cpt } \mathbb{R}^d$

$$\begin{array}{ccc}
 M & \xhookrightarrow{L} & \mathbb{R}^{d+N} \\
 \downarrow & & \\
 \text{d-mfld} & & \langle M, L \rangle \simeq \alpha_N''
 \end{array}$$

and: if α studied like as $\langle M, L \rangle$, $\exists \epsilon \frac{M}{\text{bulk}} \square$

Compute: $\pi_* \text{MO}$

Strategy: $\pi_* \text{MO}$ is an \mathbb{F}_2 -algebra.

Construct spaces:

X_N s.t. $\pi_3 X_N$ known.

it will turn out that

$$X_N = \prod_j K(\mathbb{F}_2, i_j)$$

and maps

$$\text{MO}(N) \rightarrow X_N$$

which are $H^*(-; \mathbb{F}_2)$ -isos
through a range
 $* \leq 2N$

$$\Rightarrow \text{MO}(N) \rightarrow X_N$$

π_* is through a range
 $* \leq 2N$

let $N \rightarrow \infty$

$$H^*(MO(N); \mathbb{F}_2) \cong H^*(BO(N)) \{ [v] \}$$

$$H^*(BO(N); \mathbb{F}_2) \ni w_1^{i_1} \cdots w_N^{i_N} [v]$$

as A -modules

$$\text{So } H^* MO(N) \cong (w_N) \subset H^* BO(N)$$

as A -modules

Problem: compute A -module structure of $H^* MO(k)$

Lexographical order:

$$w_k^{i_k} w_{n-1}^{i_{n-1}} \cdots w_1^{i_1} \quad S_q^I w_k = Q_I w_k$$

Claim $Q_I = w_{i_1} w_{i_2} \cdots w_{i_r} + \text{lower terms}$

$$\underline{r=1} \quad S_q^{i_1} w_k = w_{i_1} w_k \quad Q_{i_1} = w_{i_1}$$

Inductively $S_q^I w_k = S_q^{i_1} \left(\overbrace{S_q^{i_2} \cdots S_q^{i_r}}^{I'} w_k \right)$

$$= S_q^{i_1} (Q_{I'}, w_k)$$

$$Q_{I'} = w_{i_2} \cdots w_{i_r} + \cdots$$

$$S_{\mathcal{Q}}^{\alpha}(w_k Q_{I'}) = \sum_{a+b=i_1} S_{\mathcal{Q}}^a(Q_{I'}) S_{\mathcal{Q}}^b(w_k)$$

$$= \sum_{\substack{a+b=i_1 \\ i_1 \leq k}} S_{\mathcal{Q}}^a(Q_{I'}) w_b w_k$$

$$\Rightarrow Q_I = \sum_{\substack{a+b=i_1}} S_{\mathcal{Q}}^a(Q_{I'}) w_b$$

$$a=0 \rightarrow w_{i_1} \cdots w_{i_r} + \text{lower terms}$$

Claim: $S_{\mathcal{Q}}^m w_s = \sum_p w_p w_q$ w/

$$m=0 \quad \checkmark \quad p, q < 2s$$

$$S_{\mathcal{Q}}^m(w_s) = \sum_{i=0}^m \binom{m-s}{i} w_{m-i} w_{s+i}$$

$$m=s \Rightarrow w_s^s \quad \checkmark$$

$$m < s \Rightarrow s+i \leq s+m < 2s$$

$$m-i \leq m < s$$

✓

Thus $S_{\mathcal{Q}}^a Q_{I'} = S_{\mathcal{Q}}^a (w_{i_1} \cdots w_{i_r} + \text{other terms})$

*invol
w_j j < i_1*

= terms involving w_j
for $j < 2i_2$

\sum $b < i_1$

$$\underbrace{Sq^a(Q_{I'})}_{\text{index}} w_b \quad \leftarrow \begin{array}{l} \text{so these terms are inferior} \\ \text{to } v_{i_1} \dots v_{i_r} \end{array}$$

$b < i_1$

$w_i \quad j < 2i_2 \leq i_1$ \square

Thus:

$$A_{i_2} \longrightarrow (w_k) \subset H^*(BO(k))$$

$$Sq^I \longrightarrow w_{i_1} \dots w_{i_n} + \dots$$

$i \leq k$ $(\text{so } Sq^I \text{ is linearly independent})$

$$\Rightarrow MO(k) \xrightarrow{[v_n]} K(\mathbb{F}_2, k)$$

$$[u_n] \longleftrightarrow L_k$$

$$\tilde{H}^* MO(k) \longleftrightarrow \tilde{H}^*(K(\mathbb{F}_2, k))$$

$$w_{i_1} \dots w_{i_n} [v_n] + \dots \longleftrightarrow Sq^I L_k$$

L_k^2

injection for $0 \leq * \leq 2k$

$$w_n [v_n] + \dots \longleftrightarrow Sq^k L_k$$

L_k^4

$$\text{Define } H^* MO = \varprojlim \tilde{H}^{*+k}(MO(k)) \cong \mathbb{F}_2[v_1, v_2, \dots] \tilde{\wedge}$$

$$(\text{alternately: } H^m MO := \tilde{H}^{m+k} MO(k))$$

$H^* MO$ is a \mathbb{Z} -module. for $k \gg 0$
($k > m$)

$$MO(k) \wedge MO(k_0) \rightarrow MO(k+k_0)$$

$$\begin{array}{ccc} \text{gives} & H^* MO & \xrightarrow{\Delta} H^* MO \otimes H^* MO \quad \text{"coalgebra"} \\ & \gamma \mapsto \gamma \otimes \gamma + \gamma \otimes \gamma + \text{other} & \left. \begin{array}{l} \text{map of } A\text{-modules} \\ S_2^{(i)} (x_1 \otimes x_2) \end{array} \right\} \\ & \text{could view this as "revenue and costs"} & \sum_{i_1+i_2=i} S_2^{(i_1)}(x_1) \otimes S_2^{(i_2)}(x_2) \end{array}$$

Prop: $M \cong H^* MO$ is a free A -module.

[general trick for coalgebra modules over connected hopf algebras]

(PS) $A \subset A$ sub-algebra of elts of pos dg

$$N = M/\bar{A}M$$

$$M \xrightarrow{\pi} N \xleftarrow{s} \text{---} \quad s = \text{section}$$

Claim: $A \otimes N \xrightarrow{\phi} M$ is iso.

$$a \otimes n \longmapsto a s(n)$$

Note: clearly ϕ is deg.

(this works
per prop)

ϕ epi:

(easy) inject on degre.

Claim sN generates M as an A -module

$$m \in M_{\text{ic}} \quad \pi(m - s\pi(m)) = 0$$

$$\Rightarrow m - s(\pi(m)) = \sum a_j m'_j \quad |m'| < (m)$$

$$m'_j = \sum_{i \in I.M.} a'_i s(u_{ij}) \quad u_i \in N$$

$$\Rightarrow m = s(\pi(m)) + \sum a_i a'_i s(u_{ij})$$

2) ϕ monic today

$$\begin{array}{ccccccc}
 A \otimes N & \longrightarrow & A \otimes M & \longrightarrow & M & \xrightarrow{A} & M \otimes M \xrightarrow{\Delta} M \otimes N \\
 & & & & \downarrow & & \downarrow \\
 & & & & \phi & & \bar{\Delta} \\
 & & & & A-\hom & & A-\hom \\
 & & & & \uparrow & & \downarrow \\
 & & & & & & \\
 & & & & & & N \text{ has} \\
 & & & & & & \text{"formal"} \\
 & & & & & & A-\hom
 \end{array}$$

$$1 \otimes n \longmapsto s(n) \otimes z + z \otimes n$$

+ other

$$\Rightarrow a \otimes n \longmapsto a \cdot z \otimes n + \sum n_i \otimes n_i$$

$|n_i| < |n|$

$$\text{So } A \otimes N_k \longleftrightarrow A \otimes N \longrightarrow M \otimes N \rightarrow A \otimes N_k$$

$$a \otimes n \longmapsto (a \cdot \varphi) \otimes n$$

$$a \longmapsto a \cdot \varphi \quad \text{where}$$

$$\Rightarrow \text{this composit map} \Rightarrow \phi \text{ not}$$

□

A -module basis for $H^* MO$?

combinatorial exercise

Lemma as graded \mathbb{F}_2 -vector space

$$A \cong \mathbb{F}_2[\xi_1, \xi_2, \dots]$$

$$|\xi_i| = 2^i - 1$$

(Pf)
I admissible

$$i_1 i_2 \dots i_r$$

$$s_r = c_r$$

$$s_c \geq 0$$

$$s_{r-1} = c_{r-1} - 2\varepsilon_r$$

$$s_{r-2} = c_{r-2} - 2\varepsilon_{r-1}$$

⋮

$$s_1 = c_1 - 2\varepsilon_2$$

$$c_r = s_r$$

$$c_{r-1} = s_{r-1} + 2s_r$$

$$c_{r-2} = s_{r-2} + 2s_{r-1} + 4s_r$$

$$c_1 + c_2 + \dots + c_r = \underbrace{1}_{2^0} s_1 + \underbrace{(1+2)}_{2^1} s_2 + \underbrace{(1+2+4)}_{2^2-1} s_3$$

$$+ \dots + \underbrace{\left(1+2+\dots+2^{n-1}\right)}_{2^n-1} s_n$$

$$\left| S_E^I \right| = \left| \{s_1, \dots, s_r\} \right|$$

Consequence: $H^* MO(k) \cong \mathbb{F}_2[x_1, x_2, \dots]$

additively $|x_i| = i$

$$\cong \mathbb{F}_2[\xi_1, \xi_2, \dots] \otimes \mathbb{F}_2[x_i \mid i \neq 2^n-1]$$

additively

$$\cong A \otimes \mathbb{F}_2[x_i \mid i \neq 2^n - 1]$$

↑
additively

$$\Rightarrow M^* MO \cong A \otimes \mathbb{F}_2[x_i \mid i \neq 2^n - 1]$$

↑
A-dih

$$\Rightarrow M^* MO(k) \cong A \otimes \mathbb{F}_2[x_i \mid i \neq 2^n - 1] \{c_n\}$$

↓
through dim 2k

$$MO(k) \xrightarrow{\int_{M^* - \text{iso} \text{ thru } 2k} \Rightarrow \pi_* - \text{iso} \text{ thru } 2k} \prod K(\mathbb{F}_2, k+d)$$

(e_i \mid i \neq 2^n - 1)

$$|\prod x_i^{e_i}| \leq 2k$$

↑
d

$$\Rightarrow \pi_* MO(k) \stackrel{+ \text{ additively}}{\cong} \mathbb{F}_2[x_i \mid i \neq 2^n - 1] \{c_n\}$$

for * ≤ 2k

$$\Rightarrow \pi_x MO \stackrel{\text{u}}{=} F_2(x_i | i \neq 2^{-1})$$

↑
additively
Actully as a ring!

Schoefel-Whitney #'s

$$H_* MO := \varinjlim H_{*+k} MO(k)$$

$$H^* MO = (H_* MO)^*$$

$$\pi_x MO \rightarrow H_* MO$$

$$MO(k) \underset{2k}{\approx} \prod K(F_i, n_i)$$

\Rightarrow homeomorphisms in degree $\lesssim 2k$

$$\Rightarrow \pi_x MO \hookrightarrow H_* MO$$

Thus: elts of $H^k MO$

give functions on Ω

Invertibility of α :

$$x = y \in \Omega \iff \alpha(x) = \alpha(y)$$

$$\forall \alpha \in H^k MO$$



Prop: $\alpha = w_{i_1}^{e_1} \cdots w_{i_n}^{e_n} \in$

$$x = [M] \quad \text{dim } M = e_1 i_1 + \cdots + e_n i_n$$

$$\alpha(x) = \left\langle w_{i_1}(TM)^{e_1}, \dots, w_{i_n}(TM)^{e_n}; [M] \right\rangle$$

↑
↑
then are
called "character ft's"
of M

↓
fund. class

Corr:

$$M_1 \ni \text{coincident} \rightarrow M_2 \Rightarrow M_1 \text{ has}$$

Same char ft's

Use this to show that!

$$\mathcal{Q}_* \equiv \mathbb{F}_2 [[RP^i]] \mid i \neq 2^e - 1$$